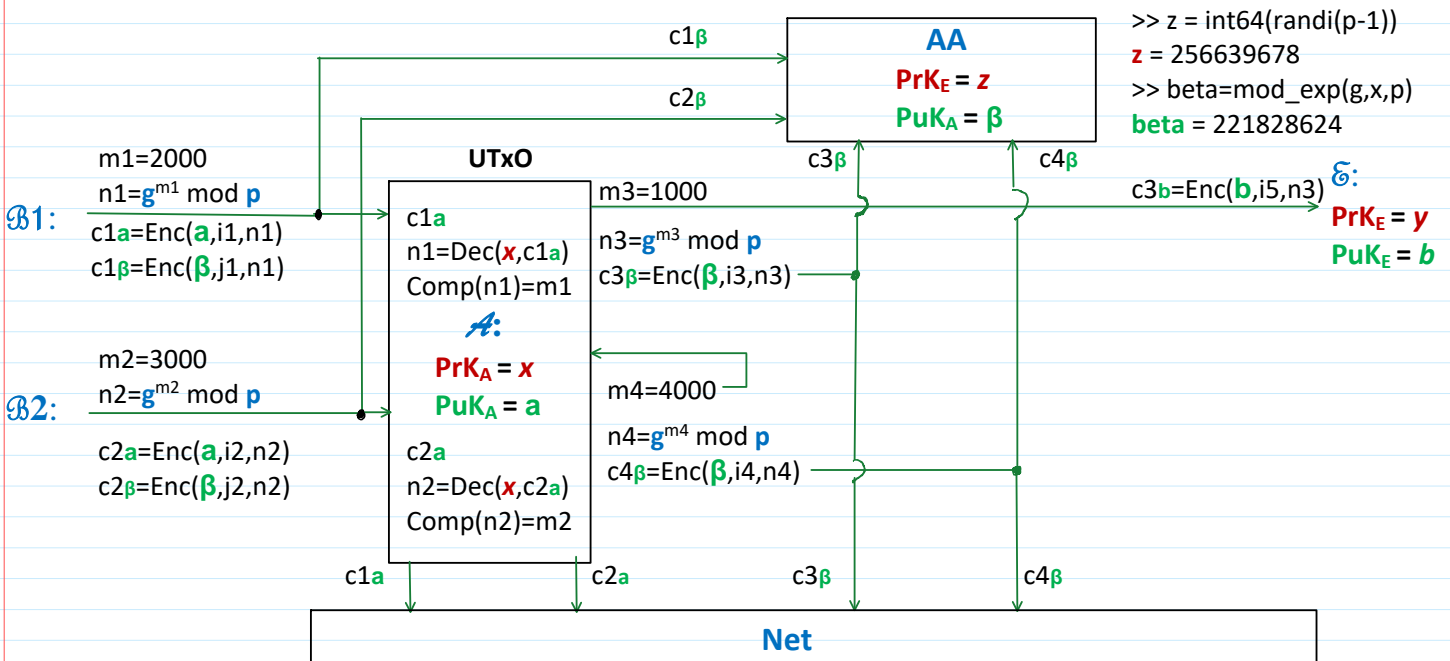


Koliokviumas: Ciphertexts equivlency proof.  
 IV-nis, Spalio 31 d., 17:30.

6027 SAKALAUŠKAS Eligijus							Pagrinc		
2024-2025 m.m. rudens semestras		9 savaitė, Spl 28		vaizdas: @savaitės / ○semestras / ○mėnuo					
Pirmadienis	Spl 28	Antradienis	Spl 29	Trečiadienis	Spl 30	Ketvirtadienis	Spl 31	Penktadienis	Lap 1
9:00									
10:30									
11:00									
12:30									
13:30	P-1706111 Kriptologija	X r.-103 P-1706127 Žuomenu sauga	XI r.-506 P-1706111 Kriptologija	XI r.-103 P-1706111 Kriptologija	XI r.-103 P-1706111 Kriptologija	XI r.-506			
15:00	prof. Eligijus SAKALAUŠKA	prof. Eligijus SAKALAUŠKA	prof. Eligijus SAKALAUŠKA	prof. Eligijus SAKALAUŠKA	prof. Eligijus SAKALAUŠKA	prof. Eligijus SAKALAUŠKA			
16:30									
17:00									
17:30		P-1706100 Kriptografinės sistemos	XI r.-103			P-1706106 r.-516 Nurodo būd žaliųjų grandinių metodu			
19:00		prof. Eligijus SAKALAUŠKA				prof. Eligijus SAKALAUŠKA			

Confidential Verifiable Transactions - 2  $PP = (p, g)$ .



$$Enc(a, i_1, n_1) = c_{1a} = (E_{1a}, D_{1a}) = (n_1 \cdot a^{i_1}, g^{i_1}) \bmod p$$

$$Enc(a, i_2, n_2) = c_{2a} = (E_{2a}, D_{2a}) = (n_2 \cdot a^{i_2}, g^{i_2}) \bmod p$$

$$c_{1a} \cdot c_{2a} = c_{12a} = Enc(a, i_{12}, n_{12}) = (E_{12a}, D_{12a}) = (n_{12} \cdot a^{i_{12}}, g^{i_{12}}) = c_{12a}$$

$$i_{12} = (i_1 + i_2) \bmod (p-1)$$

$$n_{12} = n_1 \cdot n_2 \bmod p$$

$$Enc(\beta, i_3, n_3) = c_{3\beta} = (E_{3\beta}, D_{3\beta}) = (n_3 \cdot \beta^{i_3}, g^{i_3}) \bmod p$$

$$Enc(\beta, i_4, n_4) = c_{4\beta} = (E_{4\beta}, D_{4\beta}) = (n_4 \cdot \beta^{i_4}, g^{i_4}) \bmod p$$

$$c_{3\beta} \cdot c_{4\beta} = c_{34\beta} = Enc(\beta, i_{34}, n_{34}) = (E_{34\beta}, D_{34\beta}) = (n_{34} \cdot \beta^{i_{34}}, g^{i_{34}}) = c_{34\beta}$$

$$i_{34} = (i_3 + i_4) \bmod (p-1) \checkmark$$

$$n_{34} = n_3 \cdot n_4 \text{ mod } p$$

If transaction balance is valid:  $m_1 + m_2 = 2000 + 3000 = 1000 + 4000 = m_3 + m_4$

$$\text{Then since: } n_{12} = n_1 \cdot n_2 = g^{m_1} \cdot g^{m_2} \text{ mod } p = g^{m_1+m_2} \text{ mod } p$$

$$n_{34} = n_3 \cdot n_4 = g^{m_3} \cdot g^{m_4} \text{ mod } p = g^{m_3+m_4} \text{ mod } p.$$

}  $n_{12} = n_{34} = n$

Declare Public Parameters to the network  $PP = (p, g);$

$p = 268435019; g = 2;$

$PrK_A = x \leftarrow \text{randi} \Rightarrow PuK_A = a = g^x \text{ mod } p$

$PrK_A = x \leftarrow \text{randi} \Rightarrow PuK_A = a = g^x \text{ mod } p$

```
>> p=int64(268435019)
p = 268435019
>> g=2;
```

```
>> x=int64(randi(p-1))
x = int64(220099152)
>> a=mod_exp(g,x,p)
a = 174059961
```

```
>> z=int64(randi(p-1))
z = int64(49750938)
>> beta=mod_exp(g,z,p)
beta = int64(213338364)
```

### Incomes

```
>> m1=2000;
>> n1=mod_exp(g,m1,p)
n1 = 28125784
>> i1=int64(randi(p-1))
i1 = int64(207414820)
>> a_i1=mod_exp(a,i1,p)
a_i1 = 192148999
>> E1a=mod(n1*a_i1,p)
E1a = 207347548
>> D1a=mod_exp(g,i1,p)
D1a = 202537833
```

```
>> m2=3000;
>> n2=mod_exp(g,m2,p)
n2 = 222979214
>> i2=int64(randi(p-1))
i2 = int64(67446699)
>> a_i2=mod_exp(a,i2,p)
a_i2 = 211790072
>> E2a=mod(n2*a_i2,p)
E2a = 77938423
>> D2a=mod_exp(g,i2,p)
D2a = 82080815
```

```
>> E12a=mod(E1a*E2a,p)
E12a = 52532683
>> D12a=mod(D1a*D2a,p)
D12a = 32918394
```

$C12a = (E12a, D12a)$

$c1a = (E1a, D1a)$

$c2a = (E2a, D2a)$

Verification: Dec( $x, c1a$ ) = nn1

```
>> mx=mod(-x,p-1)
mx = 48335866
>> D1a_mx=mod_exp(D1a,mx,p)
D1a_mx = 75547583
>> nn1=mod(E1a*D1a_mx,p)
nn1 = 28125784
```

Verification: Dec( $x, c2a$ ) = nn2

```
>> mx=mod(-x,p-1)
mx = 48335866
>> D2a_mx=mod_exp(D2a,mx,p)
D2a_mx = 57701660
>> nn2=mod(E2a*D2a_mx,p)
nn2 = 222979214
```

### Expenses

```
>> m3=1000;
>> n3=mod_exp(g,m3,p)
n3 = 26000000
```

```
>> m4=4000;
>> n4=mod_exp(g,m4,p)
n4 = 246627067
```

```
>> E34beta=mod(E3beta*E4beta,p)
E34beta = 57420210
>> D34beta=mod(D3beta*D4beta,p)
```

```
>> m3=1000;
>> n3=mod_exp(g,m3,p)
n3 = 260099963
>> i3=int64(randi(p-1))
i3 = int64(137379932)
>> beta_i3=mod_exp(beta,i3,p)
beta_i3 = 14259017
>> E3beta=mod(n3*beta_i3,p)
E3beta = 167897317
>> D3beta=mod_exp(g,i3,p)
D3beta = 65145889
```

```
>> m4=4000;
>> n4=mod_exp(g,m4,p)
n4 = 246637967
>> i4 = int64(randi(p-1))
i4 = int64(225960178)
>> beta_i4=mod_exp(beta,i4,p)
beta_i4 = 159771180
>> E4beta=mod(n4*beta_i4,p)
E4beta = 195130083
>> D4beta=mod_exp(g,i4,p)
D4beta = 229603826
```

```
>> E34beta=mod(E3beta*E4beta,p)
E34beta = 57420210
>> D34beta=mod(D3beta*D4beta,p)
D34beta = 107062668

C34beta = ( E34beta, D34beta)
```

```
Verification: Dec(z, c3beta) = nn3
>> mz=mod(-z,p-1)
mz = 218684080
>> D3beta_mz=mod_exp(D3beta,mz,p)
D3beta_mz = 258869169
>> nn3=mod(E3beta*D3beta_mz,p)
nn3 = 260099963
```

```
Verification: Dec(z, c3beta) = nn3
>> mz=mod(-z,p-1)
mz = 218684080
>> D4beta_mz=mod_exp(D4beta,mz,p)
D4beta_mz = 218460911
>> nn4=mod(E4beta*D4beta_mz,p)
nn4 = 246637967
```

```
>> nn12=mod(nn1*nn2,p)
nn12 = 143845522
```

$$n_{12} = n_{34} = n$$

```
>> nn34=mod(nn3*nn4,p)
nn34 = 143845522
```

$\mathcal{A}$ : must prove to the net, that  $C_{12}a$  and  $C_{34}b$  encrypted the same value  $n_{12} = n_{34} = n$ ;  $\longrightarrow$  Ciphertexts Equivalency Proof.

The statement  $st$  for this proof is the following:

$$st = \{C_{12}a, C_{34}b, a, b\}; \text{ For example: } a = g^x \text{ mod } p$$

$PubK = a$  is a statement for  $x$ .

For proof  $\mathcal{A}$  randomly generates integers  $u, v$  and  $(-v) \text{ mod } (p-1)$

$$u \leftarrow \text{randi}(\mathbb{Z}_{p-1}); \mathbb{Z}_{p-1} = \{0, 1, 2, \dots, p-2\}$$

$$v \leftarrow \text{randi}(\mathbb{Z}_{p-1})$$

$$-v \text{ mod } (p-1) \longrightarrow \Rightarrow m\bar{v} = \text{mod}(-v, p-1)$$

1. The following commitments  $\{t_1, t_2, t_3\}$  are computed:

$$t_1 = g^u \text{ mod } p$$

$$t_2 = g^v \text{ mod } p$$

$$t_3 = (D_{12}a)^u \cdot b^{-v} \text{ mod } p$$

2. The following  $h$ -value is computed using secure  $h$ -function  $H$ :

$$h = H(a || b || t_1 || t_2 || t_3)$$

```

>> u = int64(randi(p-1))      >> D12a
u = 234711265                D12a = 32918394
>> t1=mod_exp(g,u,p)         >> D12a_u=mod_exp(D12a,u,p)
t1 = 160710747              D12a_u = 190889544
>> v=int64(randi(p-1))       >> mv=mod(-v,p-1)
v = 223454508               mv = 44980510
>> t2=mod_exp(g,v,p)         >> beta_mv=mod_exp(beta,mv,p)
t2 = 131605032              beta_mv = 81562027
>> t3=mod(D12a_u*beta_mv,p)  >> t3=mod(D12a_u*beta_mv,p)
t3 = 202608126              t3 = 202608126
>> hsym = hd28('a || beta || t1 || t2 || t3')
hsym = '174059961 || 213338364 || 160710747 || 131605032 || 202608126'
hsym = 174059961 || 213338364 || 160710747 || 131605032 || 202608126
>> h=hd28(hsym)
h = 264802094

```

```

>> hsym='174059961 || 213338364 || 160710747 || 131605032 || 202608126'
hsym = 174059961 || 213338364 || 160710747 || 131605032 || 202608126
>> h=hd28(hsym)
h = 264802094

```

3.  $A$  having her  $P-K$   $x$  and  $i_{34} = (i_3 + i_4) \bmod (p-1)$  computes  $r$  and  $s$

$$r = (x \cdot h + u) \bmod (p-1)$$

$$s = (i_{34} \cdot h + v) \bmod (p-1)$$

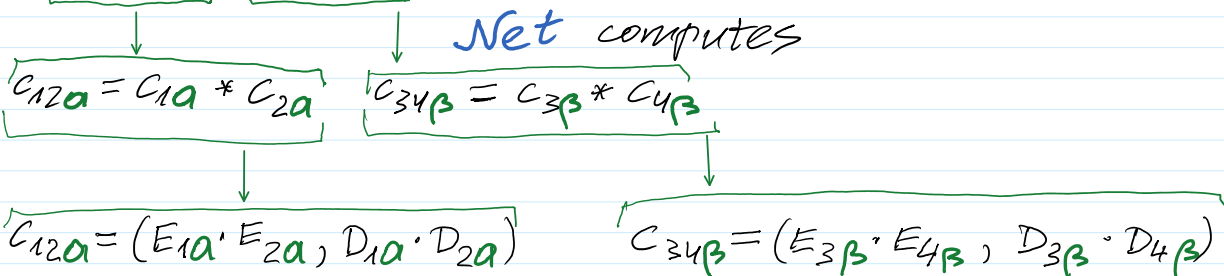
```

>> xh=mod(x*h,p-1)          >> i3
xh = 2537232                 i3 = 137379932
>> r=mod(xh+u,p-1)         >> i4
r = 237248497                i4 = 225960178
>> i34=mod(i3+i4,p-1)      >> i34h=mod(i34*h,p-1)
i34 = 94905092              i34h = 50935534
>> s=mod(i34h+v,p-1)
s = 5955024

```

and declares the following set of data to the  $Net$

$$\{C_{1a}, C_{2a}, C_{3\beta}, C_{4\beta}\} \cup \{a, \beta, t_1, t_2, t_3, r, s\} \longrightarrow Net$$



$Net$  computes  $h$ -value defined above

$$h = H(a || \beta || t_1 || t_2 || t_3)$$

Till this place

Net verifies transaction correctness by verifying the following identities

$$g^r = a^h \cdot t_1 \pmod{p}$$

$$g^s = (D_{34}\beta)^h \cdot t_2 \pmod{p}$$

$$(E_{34}\beta)^h \cdot (E_{12}a)^{-h} \cdot (D_{12}a)^r \cdot \beta^{-s} = t_3 \pmod{p}$$